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Incorporating security considerations into optimal product architecture and component sharing decision in product family design[†]

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Selecting the appropriate components to share in product family design is not a trivial decision; especially when a firm wants to protect the most sensitive information contained in their products from being exposed to users, third-party manufacturers, or undesirable agents. This article proposes tools to help designers identify sets of components containing sensitive information, as well as components that are candidates for sharing amongst the family. It also finds the most desirable component arrangement in each product that facilitates sharing while protecting the sensitive information that has been previously identified. The proposed framework is applied to three printers in which the architecture used for the ink cartridges and printheads are significantly different. Third-party manufacturers and remanufacturers offer their own alternatives for these subsystems (ink cartridges and printheads) since the customer for printer supplies is always looking for a cheaper alternative; meanwhile, the original equipment manufacturers attempt to secure their products and retain their customers with original supplies. The functional description of the system is analysed to identify the sensitive components for each printer; then, the optimal clustering strategy is found, as well as the set of components that are candidates for sharing, according to their connectivity and the security considerations.

Keywords: product family design; product security; information hiding; component sharing; product architecture

1. Introduction

Product family design is a field that traditionally encompasses the selection of a common set of components or subsystems (product platform) that can be used to develop products (variants) easily (Meyer and Lehnerd 1997) in order to fulfil a wider range of market needs. The authors explored these decisions in Rojas Arciniegas and Kim (2011), proposing a framework to obtain the set of components to share across the family, as well as the optimal architecture for each product. The work presented in this article is a logical progression from that framework including

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a new aspect of securing sensitive components into the decision of establishing the components to share and the optimal architecture for multiple products.

The architecture of a product is understood as the allocation of the functions in the physical components of the product (Ulrich 1995), and it is a key aspect for family design since it affects all the variants developed from the platform. However, the connectivity between components is not the only aspect to consider when determining the components to share in a family. In many cases, there exists sensitive information in each product that the manufacturer would not want to be accessible by the user or third-party manufacturers; then, the security of those sensitive components or interfaces becomes a concern for the designers.

These security concerns are also evident in the military domain where the term 'Anti-Tamper' is usually used to denote the requirements (or expectations) of a system or subsystem to deter any efforts to reverse engineer, exploit, or develop countermeasures for critical technology and should, in this manner, add longevity to the system (US-DoD 2010). Nevertheless, the guidelines and requirements have been developed mostly for software and electronics in which the programmable nature of these areas facilitate techniques such as encryption, code obfuscation, hardware keys, and dynamic reconfiguration (Atallah *et al.* 2004, Porter *et al.* 2009). In the physical/mechanical domain it is still a challenge to translate these protections even when the ideas and concerns are still valid.

Some manufacturing techniques allow encapsuling some components in a very secure manner; custom made integrated circuits, in-mold inserts for injection molding, and Micro Electro-Mechanical Systems (MEMS) are just a few examples. Having these capabilities at hand, it becomes important to identify which component or information is sensitive and what could be the best way to protect it. This article aims to help designers secure those components or relationships that are sensitive by using modularity to hide information, while considering component sharing and optimal architecture for multiple products.

The remaining sections of the article proceed as follows: a review of the relevant literature on this topic is presented in Section 2; the sharing decision framework is introduced in Section 3; which is followed by its application to an illustrative example in Section 4; and the article concludes with Section 5, which summarizes the proposed framework and the implications of the observed results for future work.

2. Literature review

In product architectural analysis, a few terms in common use may sometimes have confusing implications; therefore, in this article a **component** is taken to mean an element (or group of elements treated as a single part) that is a constituent piece of a product; a **module** or **cluster** refers to a group of components that ideally are highly interconnected to each other but are only weakly related to the rest of the product. A **unit** is a generic term used to refer to a component that may be viewed either as a component or as a module interface, depending on the product architecture. An **interface** refers to a component inside a module being connected to a component outside the module. Finally, a **product** means a collection of modules and components that interact to perform a common task.

A common tool to model the product is the Design Structure Matrix (DSM) as indicated in Browning (2001). The DSM was introduced by Steward (1981) and it captures the interconnections of a product, relating its components in columns and rows of a matrix. Each cell represents a relationship between components, whether it means geometrical joints, electrical connections, material flow, etc. Multiple approaches have used this representation of the product to find the optimal module definition by clustering the cells with strong relationships among components (see, for example, Yu *et al.* 2007, Wang and Antonsson 2004). This topic will be further explored in Section 2.2.

When considering a family of products, the architectural analysis is crucial since it is necessary to determine the common structure (platform) for the family. Simpson (2004) highlights that the main objective when creating a platform is to facilitate the generation of new products by having a common structure from which new variants can be developed. However, trying to capture the relationships, similarities, or differences among multiple products using multiple DSMs is not an easy task. Each DSM is different in size and content, making it very difficult to relate to other DSMs. In Rojas Arciniegas and Kim (2011), the authors used the Function-Component Matrix (FCM) from Strawbridge *et al.* (2002) to relate the components from different products and find the ones that fulfil identical functions. An FCM is a binary matrix that captures the relationship between the components of a product (columns) and the functions they fulfil (rows).

Different criteria have been used for determining *which* components to share: cost considerations (Browning and Eppinger 2002, Zacharias and Yassine 2007, Moon *et al.* 2008), Bill of Materials (BOM) (Steva *et al.* 2006), product attributes (Tucker and Kim 2008), environmental concerns (Dahmus and Gutowski 2007, Kwak *et al.* 2007, Pandey and Thurston 2008), product design variables (Khajavirad and Michalek 2008, Khajavirad *et al.* 2009), component interconnections (Rojas Arciniegas and Kim 2011), etc. The previous approach for defining which components to share is explored in Section 2.3. Nevertheless, this decision may expose certain information, components, or interfaces that are sensitive for the manufacturer. Baldwin and Clark (2000) generalized the idea of information hiding used in software engineering, proposing that the module definition is a form of abstraction in which the information that is encapsulated or contained in the module becomes *hidden* for the rest of the system. The only *visible* information from a module is the one defined by the interfaces, the connections with other components outside the module.

2.1. Security and information hiding

The security of a product is a concept difficult to quantify. Harston and Mattson (2009) introduced a set of metrics defining a 'barrier' (the difficulty of extracting information from the product) and the 'time to reverse engineer' a product, based on Ohm's law and the model of an RC circuit. Ideally, making a product secure may mean maximizing the barrier to extracting sensitive information; however, the estimation of these parameters depends on the skills and resources available to the person performing the information extraction, making them very specific and somewhat subjective.

Similarly, the term 'anti-tamper' technology has been used in the military context to denote technologies that prevent unauthorized access to specific information (see Keller 2009). An attempt to design anti-tamper circuits is discussed in Porter *et al.* (2009), in which programmable circuits are dynamically reconfigured to prevent unwanted reconstruction of the functionality of the circuit by reverse engineering techniques. A more general idea is described below.

The term 'information hiding' was introduced by Parnas (1972) in the context of software engineering, and has been used as the guideline to define the modules of a program. The idea is to decompose a system into modules in order to hide information that is likely to change, or the information that is not supposed to be accessed by the rest of the program.

By isolating or encapsulating certain information, the developer is restricting access to that information to whatever can be inferred by the interfaces (information exchanged to the rest of the system) as discussed by Hughes and Shmatikov (2004); therefore, the security of a system can be improved if the module boundaries are carefully defined to isolate the sensitive information. The concept is very general and can be used in other contexts (as suggested in Baldwin and Clark 2000), including product development, as a way to handle the complexity of a design.

Usually, the information/component/interface that is sensitive in a product is easy to identify for the designer; however, trying to protect it from the user or third-party manufacturer is not so easy. Very often, the core elements of a product are highly interconnected to the rest of the system, making it difficult to enclose completely in a module. Then it becomes crucial clearly to identify and separate the internal (hidden) resources of a module from thse that can be accessed by the rest of the system or users (O'Hearn *et al.* 2009).

A simple case is discussed in Legrand (2007), in which the tradeoff between modularity and security is explored for 'protected distribution systems'. These systems are used for the protection of networks and are custom made, where applying a modular design for the conduit leads to the exposure of the network at the module boundaries, thus compromising the security of the entire system. Then, the challenge for the product designers is to find an optimal component arrangement that preserves the desired protection for those sensitive elements.

2.2. Optimal clustering for product architecture

In Wang and Antonsson (2004) and Yu *et al.* (2007), the optimal architecture for a product is achieved through the use of the Minimum Description Length (MDL) principle. This principle promotes the simplest description for any structure; thence, each component arrangement of a product can be coded by a model in which the simpler the model is, the shorter the description is. In the end, the idea is to measure the length of the description and find the minimum, which corresponds to the simplest (optimal) component arrangement for the product.

Both approaches use this same principle, but in different manners. Yu *et al.* (2007) used a description of the overall structure and accounted for two types of mismatch between a model and the actual product: type I corresponds to the relationships between components that are left outside the modules defined by the model; type II corresponds to the relationships inside the modules that are weak or do not exist (recall that components inside the modules should be strongly interconnected). The formulation is as follows:

$$f_{\text{MDL}} = (1 - \alpha - \beta) \left(n_{\text{cl}} \log_2 n_{\text{c}} + \log_2 n_{\text{c}} \sum_{i=1}^{n_{\text{cl}}} cl_i \right) + \alpha [S_1 (2 \log_2 n_{\text{c}} + 1)] + \beta [S_2 (2 \log_2 n_{\text{c}} + 1)].$$
(1)

In this formulation, the first term corresponds to the length of the overall model, while the second and third terms correspond to the mismatch lengths of types I and II, respectively. n_{cl} is the number of clusters, n_c is the total number of components of the product, cl_i is the number of elements in cluster *i*, and α and β are weighting factors for the mismatches. The terms S_1 and S_2 account for the mismatches and are defined as follows:

$$S_{1} = \sum_{d'_{ij}=1} (1 - p_{ij}); \quad S_{2} = \sum_{d'_{ij}=0} (p_{ij});$$

$$p_{ij} = \frac{d_{ij} - d_{\min}}{d_{\max} - d_{\min}}; \quad d_{\max} = \max_{ij} d_{ij}; \quad d_{\min} = \min_{ij} d_{ij};$$
(2)

where d_{ij} is the entry of the *i*th row and the *j*th column of the DSM, and p_{ij} is the same entry normalized. The entry of the ideal binary DSM described by the model is d'_{ij} , in which there are no mismatches, *i.e.* all the entries are '1' inside the clusters and '0' outside the clusters.

On the other hand, Wang and Antonsson (2004) used a formulation based on the number of links from each component. This formulation allows a representation of each component (unit) in the product as follows:

$$MDL_{j}^{(u)} = -\log_{2}\left(N_{j}^{(u)} / \sum_{k=1}^{N^{(u)}} N_{k}\right)$$
(3)

in which $N_j^{(u)}$ is the number of units connected to unit j, and $N^{(u)}$ is the number of units at the level in which the unit j is. In this context, two units are considered to be at the same level if they are inside the same module or outside every module. Since a unit can be seen as a component if it is considered inside a module, but also as an interface if it has connections outside the module, the total representation of unit j in product k would be

$$MDL_{j}^{k} = MDL_{j}^{(c)} + MDL_{j}^{(o)} = -\sum_{u=\{c,o\}} \log_{2} \left(N_{j}^{(u)} / \sum_{k=1}^{N^{(u)}} N_{k} \right),$$
(4)

where 'c' represents components and 'o' interfaces. This abstraction of the representation becomes very useful when analysing how tightly connected each element of the product is.

2.3. Optimal sharing decision making

Many criteria have been used to define the set of components to share; however, when considering product architecture, in Rojas Arciniegas and Kim (2011), the authors used the Impact Metric (IM) and the FCM to establish the suitable candidates for sharing. The IM introduced by Rojas and Esterman (2008), gives a proxy for the easiness to change a component under a given architecture. The formulation is as follows:

$$IM_{i}^{k} = MDL_{i}^{k}CI_{i}^{k},$$
(5)

in which the CI_j^k corresponds to the Coupling Index (CI) of component *j* in product *k*. The CI was introduced by Martin and Ishii (2002) and the formulation is based in the DSM that contains the sensitivity analysis for the impact of a change in the specifications of the components. This assumes that the components in the columns supply the specification and the components in the rows receive the specifications. Each row and column of the DSM is added giving the CI-Receiving (CIR) for the rows, and the CI-Supplying (CIS) for the columns. The total CI for each components is the sum of the CIS and the CIR as follows:

$$CI_{j}^{k} = CIS_{j}^{k} + CIR_{j}^{k} = \sum_{i=1}^{n_{c}^{k}} DSM^{k}(i, j) + \sum_{i=1}^{n_{c}^{k}} DSM^{k}(j, i).$$
(6)

The IM for each component of each product is an indication of how easy it is to change a particular component. Nevertheless, in order to identify common components across different products, it is necessary to find common ground. Therefore, the authors used the FCM to match components from different products by the functions they fulfil. The criterion was to find the set of components that fulfil the same functions in a pair of products, and analyse if the sum of IM scores from the components involved was below a threshold set by the designer. In this manner, the pairwise evaluation was conducted for all the products in the family. In the end, the set of components to share from all the products was established and stored in a binary decision vector (Y^k) for each product k.

In Rojas Arciniegas and Kim (2011), after the decision vectors were established, an optimization problem was solved for each product of the family with a dual objective: on one hand, finding

an optimal clustering of the components, making use of the product MDL representation of Equation (1); meanwhile, the decision vector was multiplied by the IM vector from Equation (5) selecting the components to include in the objective function, allowing those components to be changed easily in the product. The optimization problem was solved using genetic algorithms and the chromosome was a binary string that represented a particular component arrangement in clusters for product *k*, also represented in a binary matrix form (X^k) . In this matrix, the columns are the components and the rows are the clusters; therefore, each cell (i, j) indicates if the component *j* belong to the cluster *i*. The objective function was defined as follows:

$$f(X^k) = f_{\text{MDL}}^k + (Y^k)^{\text{T}} \text{IM}^k.$$
(7)

Since the IM depends on the current architecture of the product, the process was iterative. At each iteration (*t*) a new component arrangement was achieved for each product, progressively establishing the final list of components to share. The process was repeated until the change in the decision vectors for all the products being considered was below a given tolerance ϵ :

$$\sum_{k=1}^{m} \frac{\|Y^{k}(t) - Y^{k}(t-1)\|^{2}}{n_{c}^{k}} \le \epsilon.$$
(8)

3. Generating security constraints for optimal product family design

The framework presented in this section builds upon the work presented in Rojas Arciniegas and Kim (2011) (also described in Section 2.3). It incorporates security considerations when selecting the optimal candidates for sharing across multiple products, simultaneously identifying the optimal component arrangement for each individual product.

3.1. Defining sensitive components

Rojas Arciniegas and Kim (2011) proposed the use of the functional mapping captured in the FCMs to automatically identify the components that could be matched between products based on the functions they perform. Since the FCM for each product is already available (or at least it is a requirement for the application of the framework), this information can be used to designate the sets of sensitive components, by relating the functions that are 'critical' to the components that perform those in each product. Thus, components can be grouped in sets based on the same functional mapping.

The main difference with the approach presented in Rojas and Kim (2010) is that the sets of sensitive components are automatically generated using the FCM and designating the 'critical' functions, which are independent of the realization of each product. The term *critical* is used to denote the functions that carry most of the intellectual property of the product and therefore should be protected. These functions are easier to identify than individual components in each product, since those are the functions in which companies usually invest more resources and have heavy research and development.

The critical functions are designated by means of a binary vector (*CrFun*) for which a one (1) would indicate that the function is critical. The size of the vector is equal to the number of functions (n_f) considered for the analysis (also the number of rows in the FCMs). *CrFun* is the input required from the designer and corresponds to the nature of the company and where the intellectual property is.

The automated detection of the sensitive components and its clustering in restricted sets is performed (similarly to the functional matching proposed in Rojas Arciniegas and Kim 2011) by examining the vector *CrFun* for critical functions in the following manner.

- A vector keeps track of the critical functions evaluated throughout the selection of the restricted sets, denoted as the *master function-vector* (*f-v*).
- Another vector marks the critical functions related to the matching components, denoted as the *local f-v*, to form each restricted set.
- A different vector is used to register the components related to the critical functions in the *local f*-*v* forming each restricted set. This vector is denoted as the *component-vector* (c-v).
- A matrix is formed row by row with the results from the *c*-*v* and it is returned at the end of the identification process, in which the rows indicate the restricted sets (*n*_{rcl}) of components for product *k*. This matrix is denoted as the *security matrix* (*SecM*^k).

The master f-v and the local f-v are binary vectors with n_f elements, while the c-v is a vector with n_c elements. The security matrix (*SecM*^k), is a binary matrix of dimensions n_{rcl} by n_c . An example of a *SecM* for a product with eight components is shown in Table 1, in which components B, C and D are required to be inside a module. Similarly, components F, G and H are also required to be in a module. The requirement does not imply that the modules should be different from one another – it could be one module that contains both sets, but the sets have to be fully contained in the module.

The selection of restricted sets of components is conducted as follows.

- Step 1. Initialize the *master f-v* with a zero vector and start by evaluating the first function (i = 1).
- Step 2. Check if the function is critical; if yes, continue; if not, go to Step 14.
- Step 3. Check if the function has been evaluated before (*master f-v(i)=1?*); if yes, go to Step 14; if not, continue.
- Step 4. Mark the *master f-v* with a '1' for the corresponding function, and initialize the *local f-v* and *c-v* with zero vectors.
- Step 5. Mark the *local f-v* with a '1' for the corresponding function, and set t = i.
- Step 6. Check if the function should be considered (local f v(t) = 1?); if yes, continue; if not, go to Step 11.
- Step 7. Find components related to the function following the row of the FCM corresponding to the function being evaluated until a '1' is found.
- Step 8. Mark the corresponding component in the c-v.
- Step 9. Find other critical functions related to the component following the column of the corresponding component and mark them in the *local f-v*.
- Step 10. Return to Step 7 until all the components have been evaluated.
- Step 11. Continue to the next function (t = t + 1) and return to Step 6 until all the functions in the *local f-v* have been evaluated.
- Step 12. Mark the functions marked in the local f-v into the master f-v.
- Step 13. Check if the c-v is empty; if not, create a new row for $SecM^k$ and mark the components in the c-v into the row of the matrix.

				Comp	onents			
	A	В	С	D	Е	F	G	Н
Set1	0	1	1	1	0	0	0	0
Set2	0	0	0	0	0	1	1	1

Table 1. Example of the security matrix (*SecM*) for a product of eight components.

Step 14. Continue the evaluation of the next function (i = i + 1) and return to Step 2 until all the functions have been evaluated.

In order to achieve this restriction, a new term is added in the MDL formulation. This term accounts for the mismatch of the current architecture to the constraint designated in the *SecM* for the product. The formulation in Equation (1) is redefined as follows:

$$\hat{f}_{\text{MDL}} = (1 - \alpha - \beta - \gamma) \left(n_{\text{cl}} \log_2 n_{\text{c}} + \log_2 n_{\text{c}} \sum_{i=1}^{n_{\text{cl}}} cl_i \right) + \alpha [S_1 (2 \log_2 n_{\text{c}} + 1)] + \beta [S_2 (2 \log_2 n_{\text{c}} + 1)] + \gamma [S_3 (\log_2 n_{\text{c}} + \log_2 n_{\text{rcl}} + \log_2 n_{\text{cl}})],$$
(9)

in which γ is a weighting factor, $n_{\rm rcl}$ is the number of restricted sets of components (rows of *SecM*), and S_3 accounts for the mismatches with the restricted sets. The description of this mismatch would be given by the component number, the restricted set number, and the cluster number; therefore, the description length is given by $(\log_2 n_c + \log_2 n_{\rm rcl} + \log_2 n_{\rm cl})$. S_3 counts the number of these mismatches in the following manner:

$$S_{3} = \sum_{i=1}^{n_{\text{cl}}} \sum_{j=1}^{n_{\text{cl}}} \left[SecM_{i.}(e - X_{j.})^{\text{T}} \min \left\{ \min \left\{ SecM_{i.}X_{j.}^{\text{T}}, SecM_{i.}(e - X_{j.})^{\text{T}} \right\}, 1 \right\} \right] + \sum_{i=1}^{n_{\text{cl}}} \left[SecM_{i.}e^{\text{T}} \left(1 - \min \left\{ \sum_{j=1}^{n_{\text{cl}}} SecM_{i.}X_{j.}^{\text{T}}, 1 \right\} \right) \right],$$
(10)

where $SecM_i$. corresponds to the *i*th row vector of SecM, X_j . the *j*th row vector of X, and *e* is a unity row vector of the appropriate dimension. The first term accounts for the mismatches between each cluster (X_j) and each restricted set $(SecM_i)$, preventing overcounting mismatches when no component in the restricted set is in the cluster. The second term penalizes the restricted sets that are not included in any of the clusters.

This modification to the original MDL formulation finds an optimal clustering for the product while maintaining the constraints defined in *SecM* through the optimization process.

3.2. Component sharing

The component sharing decision is also affected by the restriction on the sensitive components. In Rojas Arciniegas and Kim (2011), a candidate for sharing gets isolated in a cluster by itself to facilitate the sharing process. Having a sensitive component selected as a candidate would violate the restriction imposed by *SecM*, unless it is selected along with all the other components of the restricted set and it would have to be shared as a module. Therefore, it is necessary to redefine the decision vector Y^k , so this aspect is taken into consideration. The reformulation is as follows:

$$Y_{c}^{k}(j) = Y^{k}(j) \prod_{i=1}^{n_{rcl}} \left(1 - SecM^{k}(i, j)\right)$$

$$Y_{m}^{k}(l) = \min\left\{\sum_{i=1}^{n_{rcl}} \left[\left(1 - \min\left\{\left(SecM^{k}(i, \cdot)e - X^{k}(l, \cdot)SecM^{k}(i, \cdot)'\right), 1\right\}\right) \times \left(1 - \min\left\{\left(SecM^{k}(i, \cdot)e - Y^{k}SecM^{k}(i, \cdot)'\right), 1\right\}\right)\right], 1\right\}$$

$$\forall i \in \{1, \dots, n_{rcl}\}, \quad j \in \{1, \dots, n_{c}\}, \quad l \in \{1, \dots, n_{cl}\},$$
(11)

where Y_c^k exclude all the restricted components of product k from Y^k ; Y_m^k indicates which modules contain restricted sets entirely and are suitable candidates according to Y^k . The original decision vector Y^k is still obtained by pairwise comparisons between products, establishing a functional matching of components from different products and evaluating if the sum of the IM score for the components related is below a given threshold. This analysis determines the components that are easy to change under the current architecture, making those candidates for sharing across the family of products. (Refer to Section 3.3 of Rojas Arciniegas and Kim 2011 for more details.)

The IM is also computed for the modules defined in X^k by treating each cluster as a component, aggregating the external connections, computing the CI for the cluster, and multiplying it by the MDL of the module. The result of this calculation is stored in the vector IM_m^k .

The subproblem (clustering optimization for each product k) is redefined as follows:

$$\min f^k = \alpha_1 Y_c^k \mathbf{I} \mathbf{M}^k + \alpha_2 Y_m^k \mathbf{I} \mathbf{M}_m^k + \hat{f}_{MDL}^k, \tag{12}$$

where α_1 and α_2 are weighting factors. This objective function determines the optimal clustering strategy for the individual products based on the DSM of the product (fixed), the restrictions for sensitive components defined in the *SecM* (fixed), the modified decision vector passed from the master problem (component sharing) as a parameter, and evaluates the chromosome (X^k) during the process of optimization until finding the optimal clustering strategy.

3.3. Framework overview

The framework defined in Rojas Arciniegas and Kim (2011) requires adjustments to accommodate the new criteria involved in the analysis. A new matrix *SecM* is required for each product, and the objective function in the optimization process requires two important modifications to include the security considerations: first, to incorporate a mismatch between the restricted sets in *SecM* and the clustering strategy X for the MDL model; and second, to alter the decision vector to exclude the sensitive components that are selected for sharing without all the elements in the corresponding restricted set. Therefore, the overall framework proceeds as follows (Figure 1).

- Step 1. Construct the product DSM for each variant of the family.
- Step 2. Construct the FCM for each product and include all the functions in the family.
- Step 3. Define the set of critical functions in the binary vector CrFun.
- Step 4. Generate the restricted sets of components (SecM) for each product.
- Step 5. Initialize the modified decision vectors Y_c^k and Y_m^k with zero vectors.
- Step 6. Run the optimization for each product separately. Include both the modified MDL representation of the product (\hat{f}_{MDL}^k) and the IM of the modified selection of components and modules for sharing $(\alpha_1 Y_c^k IM^k + \alpha_2 Y_m^k IM_m^k)$, following Equation (12) to obtain a new clustering strategy X^k that considers the security of the sensitive components and component sharing.
- Step 7. Run the component sharing selection algorithm to obtain the decision vectors Y for each product in the analysis. This algorithm calculates the IM for all the components and performs functional matching in order to find candidates for sharing with low IM. The selection of the candidates depends on a threshold level for the IM.
- Step 8. Obtain the modified decision vectors Y_c^k and Y_m^k for each product.
- Step 9. Verify if a stopping criterion is met (*i.e.* change in all the decision vectors below tolerance ϵ); if not, return to Step 6.



Figure 1. Framework overview flow diagram.

Ultimately, the proposed framework provides designers with an optimal component arrangement for each product and a set of candidates for components to share across multiple products in considerations of security (sensitive components), product architecture (internal connectivity), and component sharing (family of products).

4. Illustrative example

The framework discussed in the previous section is applied to analysing HP inkjet cartridges and their architecture for three distinct printers: the HP DeskJet 1220C; the HP Photosmart Plus; and the HP Business Inkjet 2600 (see Figure 2).

These three printers exemplify the three strategies that HP currently uses for inkjet cartridge design: the first uses black (BK) and tricolour – cyan (C), magenta (M) and yellow (Y) – cartridges that include the printhead; the second uses individual ink cartridges for each colour but the printhead is fixed to the printer; the last printer uses individual ink cartridges as well as printheads that are replaceable separately. The supplies for each printer are shown in Figure 2, as well as diagrams that show the ink cartridges, printheads and carriages. This figure also shows the direction in which the compartment moves and the elements that are carried by it, and how the printing media pass by the printheads. These diagrams do not include the body of the printer or any subsystem other than the elements mentioned above, clarifying the difference in architecture of the three printers.

Each architecture is significantly different and includes distinctive features depending on the intended market for that product. The Deskjet printer is intended for basic colour printing for a user at home or in an office; the Photosmart is directed to users who require to print photos as well as documents, but not at a professional level; lastly, the business inkjet is intended for higher volumes of colour printing in an office.

In a recent market report, Gartner Inc. affirms that printer OEMs (Original Equipment Manufacturers) could lose more than \$13 billion to third-party manufacturers of supplies between mid-2009 and mid-2010 (Gartner 2009). This potential loss puts in perspective the importance for OEMs to secure their products and prevent third-party manufacturers or remanufacturers copying their intellectual property.



Figure 2. Printers and cartridge supplies for each printer: (a) Printer HP Deskjet 1220C and cartridges HP45 (BK) and 78 (C, M, Y); (b) Printer HP Photosmart Plus and cartridges HP564 (BK, Photo BK, C, M, Y); (c) Printer HP Business Inkjet 2600 and cartridges HP10 (BK) and 11 (C, M, Y), Printheads HP11 (BK, C, M, Y).

4.1. Application of the framework

- Steps 1 & 2. The DSMs and FCMs were developed, only for the subsystems related to the ink cartridges and printheads, for the three printers mentioned above. The matrices were developed based on the service manuals and customer manuals for the printers (see HP 2000, HP 2009 and HP 2002, respectively). The DSMs and FCMs are not shown owing to space limitations.
 - Step 3. The critical functions were defined in *CrFun*, indicating the functions that are most sensitive for the printing subsystems, independent of the actual implementation in each printer. The selection of the critical functions is explained further in Section 4.2 and the *CrFun* vector can be seen in Table 2.
 - Step 4. The security matrices (*SecM*) were generated for each printer, reflecting the most sensitive components in the cartridges and the architecture of each printer. The *SecM*'s can be seen in Tables 4, 5 and 6, along with a summary of the numerical results for each printer.
 - Step 5. The modified decision vectors $(Y_c^k \text{ and } Y_m^k)$ were initialized with zero vectors of the appropriate dimension for each product.
 - Step 6. The subproblems were solved using genetic algorithms with the modified decision vector for each printer, the product DSM, and the *SecM* as parameters for the

Function #	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	2
0.5	Store Kink	Store C ink	Store M ink	Store Y ink	Move carriage along paper	Deliver K Ink precisely into paper	Deliver C Ink precisely into paper	Deliver M Ink precisely into paper	Deliver Y Ink precisely into paper	Send/Receive signals	Provide support for movement	Provide support for cartridges	Protect K ink from environment	Protect C ink from environment	Protect M ink from environment	Protect Y ink from environment	Protect K printhead from environment	Protect C printhead from environment	Protect M printhead from environment	Protect Y printhead from environment	Clean printheads	Prevent ink spills	Absorb extra ink	Sense K ink level and cartridge info	Sense C ink level and cartridge info	Sense M ink level and cartridge info	Sense Y ink level and cartridge info	Transnort ink from reconneir to printhes
CrFun	1	1	1	1	0	1	1	1	1	0	0	0	1	1	1	1	1	1	1	1	0	0	0	1	1	1	1	1

Table 2. Critical functions for the printing subsystem.

optimization. The optimal clustering strategy for each product was returned as the optimal chromosome (X^k) according to Equation (12), which captures connectivity, component sharing and security restrictions.

- Step 7. The optimal chromosomes from all products were passed to the master problem as parameters and the component sharing selection algorithm was executed. The threshold value for the IM was set at 1.0 for this example. The decision vectors (Y^k) were returned for each product and can be seen in Tables 4, 5 and 6.
- Step 8. The modified decision vectors $(Y_c^k \text{ and } Y_m^k)$ were obtained according to Equation (11).
- Step 9. The tolerance for the stopping criteria, in this case, was selected to be zero ($\epsilon = 0$). Having this tolerance, Steps 6 to 8 were repeated for two iterations, after which there was no change in any decision vector.

4.2. Identifying the critical functions

In order to establish the security restrictions in all three printers it was necessary to identify the critical functions of the printing subsystem. First, the functional structure was examined designating the flow of information, energy and material through the printing system. The list of functions came from the FCMs since that is a set of functions considered for all printers and they are independent of the physical implementation of the individual products.

In Figure 3 the functional structure reflects how the functions involved in handling the ink and taking each colour from the reservoirs to the paper in a precise manner were the most critical functions in this subsystem. Therefore, these functions are designated as critical functions in the vector *CrFun* (see Table 2), which is later used to generate the security matrices for each printer (*SecM*^k).

Table 3 shows how the critical functions are related to the components of the printer HP Photosmart Plus, generating the security matrix ($SecM^2$), which is composed of five restricted sets, the first related to the printheads and the remaining four related to the individual ink cartridges. The generation of this *SecM* was as follows.

Step 1. i = 0 and master $f - v = \hat{0}$ ($\hat{0}$ is the zero vector of n_f elements).

- Step 2. Is function *i* critical? CrFun(1) = 1.
- Step 3. Is master f-v(i) = 1? master f-v(1) = 0, then continue.
- Step 4. *master* f-v(1) = 1, *local* f- $v = \hat{0}$, and c- $v = \hat{0}$.
- Step 5. local f v(1) = 1, t = i = 1.
- Step 6. Is $local f \cdot v(t) = 1$? $local f \cdot v(1) = 1$, then continue.

	Se	ecM	12			28	27	26	25	24	23	22	21	20	19	18	17	16	5	14	13	12	11	10	9	00	~	6	ъ	4	ω	2		Function #	
Set 5	Set 4	Set 3	Set 2	Set 1		Transport ink from reservoir to printhead 1	Sense Y ink level and cartridge info 1	Sense M ink level and cartridge info 1	Sense C ink level and cartridge info 1	Sense K ink level and cartridge info 1	Absorbe extra ink 0	Prevent ink spills 0	Clean printheads 0	Protect Y printhead from environment 1	Protect M printhead from environment 1	Protect C printhead from environment 1	Protect K printhead from environment 1	Protect Y ink from environment 1	Protect M ink from environment 1	Protect C ink from environment 1	Protect K ink from environment 1	Provide support for cartridges 0	Provide support for movement 0	Send/Receive signals 0	Deliver Y Ink precisely into paper 1	Deliver M Ink precisely into paper 1	Deliver C Ink precisely into paper 1	Deliver K Ink precisely into paper 1	Move carriage along paper 0	Store Y ink 1	Store M ink 1	Store C ink 1	Store K ink 1	CrFun	Component #
0	0	0	0	0																		×	×				Ĩ							Carriage	1
0	0	0	0	1										×	×	×	×					×							Ĵ					Plastic case Ph	N
0	0	0	0	1													×											×	1					Printhead K	ω
0	0	0	0	1					Γ					×	×	×									×	×	×							Printhead CMY	4
	0	0	0	0					T																									Secure Printheads Lever	5
0	0	0		0						×											×												×	Ink depository K	6
0	0	0	1	0	1																×													Plastic case K	2
0	0	0	н	0	1					×																								Smart chip K	00
0	0	4	0	0					×											×												×		Ink depository C	9
0	0	-	0	0										1						×														Plastic case C	10
0	0	-	0	0					×																									Smart chip C	11
0	-	0	0	0				×											×										Ì		×			Ink depository M	12
0	E	0	0	0															×															Plastic case M	13
0	Ţ	0	0	0				×																										Smart chip M	14
1	0	0	0	0			×											×												×				Ink depository Y	15
1	0	0	0	0														×											ĺ					Plastic case Y	16
1	0	0	0	0			×																											Smart chip Y	17
0	0	0	0	0																				×										Carriage PCB	18
0	0	0	0	0																			×											PCB Holder	19
0	0	0	0	0																									×					Belt attach	20
0	0	0	0	0																				×										Flex cable	21
0	0	0	0	0																									×					Carriage Belt	22
0	0	0	0	0																			×											Slider	23
0	0	0	0	0																									×					Carriage Motor	24
0	0	0	0	0								×																						Serv. Station Body	25
0	0	0	0	0								×		_																				Top Case	26
0	0	0	0	0									×																					Wiper	27
0	0	0	0	0							×																							Sponge	28
0	0	0	0	0									×																					Pallet	29
0	0	0	0	0									×																					Shaft	щ
0	0	0	0	0									×																					Motor & Gears	31

Table 3. FCM, CrFun and SecM generation for the printer HP Photosmart Plus.

Engineering Optimization

Step 7. Follow row 1 of FCM until a '1' is found, FCM(1, 6) = 1.

Step 8. c - v(6) = 1.

Step 9. Other related functions of component 6 which are also critical: functions 13 and 24, then $local f \cdot v(13) = 1$, $local f \cdot v(24) = 1$.

 Table 4. Results summary for the printer HP Photosmart Plus.

				X2	(2)									Y2			Se	ecN	12			
Cluster 10	Cluster 9	Cluster 8	Cluster 7	Cluster 6	Cluster 5	Cluster 4	Cluster 3	Cluster 2	Cluster 1		Yc2(2)		Y2(2)	Y2(1)	Y2(0)	Set 5	Set 4	Set 3	Set 2	Set 1		Comp. #
0	0	0	0	0	0	0	0	0	0		0		0	0	0	0	0	0	0	0	Carriage	1
0	0	0	0	0	0	0	0	0	1		0		0	0	0	0	0	0	0	1	Plastic case Ph	2
0	0	0	0	0	0	0	0	0	1		0		0	0	0	0	0	0	0	1	Printhead K	ω
0	0	0	0	0	0	0	0	0	1		0		0	0	0	0	0	0	0	1	Printhead CMY	4
0	0	0	0	0	0	0	0	0	0		0		0	0	0	0	0	0	0	0	Secure Phd Lever	ъ
0	0	0	0	0	0	0	0	1	0		0		1	1	0	0	0	0	1	0	Ink depository K	6
0	0	0	0	0	0	0	0	1	0		0		1	1	0	0	0	0	1	0	Plastic case K	7
0	0	0	0	0	0	0	0	1	0		0		1	1	0	0	0	0	1	0	Smart chip K	∞
0	0	0	0	0	0	0	1	0	0		0		1	1	0	0	0	1	0	0	Ink depository C	9
0	0	0	0	0	0	0	1	0	0		0		1	1	0	0	0	1	0	0	Plastic case C	10
0	0	0	0	0	0	0	1	0	0		0		-	1	0	0	0	1	0	0	Smart chip C	11
0	0	0	0	0	0	1	0	0	0		0		1	1	0	0	1	0	0	0	Ink depository M	12
0	0	0	0	0	0	1	0	0	0		0		1	1	0	0	1	0	0	0	Plastic case M	13
0	0	0	0	0	0	1	0	0	0		0		1	1	0	0	1	0	0	0	Smart chip M	14
0	0	0	0	0	1	0	0	0	0		0		1	1	0	1	0	0	0	0	Ink depository Y	15
0	0	0	0	0	1	0	0	0	0		0		1	1	0	1	0	0	0	0	Plastic case Y	16
0	0	0	0	0	1	0	0	0	0		0		1	1	0	1	0	0	0	0	Smart chip Y	17
0	0	0	0	0	0	0	0	0	0		0		0	0	0	0	0	0	0	0	Carriage PCB	18
0	0	0	0	0	0	0	0	0	0		0		0	0	0	0	0	0	0	0	PCB Holder	19
0	0	0	0	1	0	0	0	0	0		1		1	1	0	0	0	0	0	0	Belt attach	20
0	0	0	0	0	0	0	0	0	0		0		0	0	0	0	0	0	0	0	Flex cable	21
0	0	0	1	0	0	0	0	0	0		1		1	1	0	0	0	0	0	0	Carriage Belt	22
0	0	0	0	0	0	0	0	0	0		0		0	0	0	0	0	0	0	0	Slider	23
0	0	1	0	0	0	0	0	0	0		1		1	1	0	0	0	0	0	0	Carriage Motor	24
0	1	0	0	0	0	0	0	0	0		0		0	0	0	0	0	0	0	0	Serv. Station Body	25
0	1	0	0	0	0	0	0	0	0		0		0	0	0	0	0	0	0	0	Top Case	26
0	1	0	0	0	0	0	0	0	0		0		0	0	0	0	0	0	0	0	Wiper	27
÷	0	0	0	0	0	0	0	0	0		1		нi)	1	0	0	0	0	0	0	Sponge	28
0	0	0	0	0	0	0	0	0	0		0		0	0	0	0	0	0	0	0	Pallet	29
0	0	0	0	0	0	0	0	0	0		0		0	0	0	0	0	0	0	0	Shaft	30
0	0	0	0	0	0	0	0	0	0		0		0	0	0	0	0	0	0	0	Motor & Gears	31
0	0	0	0	0	1	1	1	1	0	Ym	2(2	!)	1									

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	Comp. #	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	
		Carriage	plastic case K	ink depository K	printhead K	plastic case CMY	ink depository CMY	printhead CMY	ESD Spring	ESD Blade	Carriage PCB	PCB Holder	Slider attach	Belt attach	Flex cable	Slider	Carriage Motor	Carriage Belt	Serv. Station Body	Top Case	Wiper	Sponge	Pallet	Shaft	Motor & Gears	
N	Set 1	0	1	1	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
Se	Set 2	0	0	0	0	1	1	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
												_		_												
	Y1(0)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
Y1	Y1(1)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0	
	Y1(2)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0	
	Yc1(2)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0	1(2
																										۳,
	Cluster 1	0	1	1	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
(2)	Cluster 2	0	0	0	0	1	1	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
X1	Cluster 3	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	0	0	0	0	0
	Cluster 4	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0	0

- Step 10. Return to Step 7 to continue to follow row 1 looking for a '1'; since no other 1 is found, continue to Step 11.
- Step 11. Increment *t* and return to Step 6 looking for *local f*-v(t) = 1.
- Step 6. At t = 13, $local f \cdot v(13) = 1$, then continue.
- Step 7. Follow row 13 of FCM until a '1' is found, FCM(13, 7) = 1.
- Step 8. c v(7) = 1.
- Step 9. Other related functions to component 7: no other.
- Step 10. Return to Step 7 to look for other components related to function 13, since no other function is found continue to Step 11.
- Step 11. Increment t and return to Step 6 looking for local f-v(t) = 1.
- Step 6. At t = 24, $local f \cdot v(24) = 1$, then continue.
- Step 7. Follow row 24 of FCM until a '1' is found, FCM(24, 8) = 1.
- Step 8. c v(8) = 1.
- Step 9. Other related functions to component 8: no other.
- Step 10. Return to Step 7 to look for other components related to function 24, since no other function is found continue to Step 11.
- Step 11. Increment t and return to Step 6 looking for $local f \cdot v(t) = 1$, since no other t makes $local f \cdot v(t) = 1$, then continue to Step 12.
- Step 12. master f-v(13) = 1, master f-v(24) = 1.
- Step 13. Since the *c*-*v* is not empty then the *c*-*v* is copied into the first row of SecM, SecM(1, 6) = 1, SecM(1, 7) = 1 and SecM(1, 8) = 1.
- Step 14. Increment *i* and return to Step 2 looking for critical functions to generate other secured sets.

Following the same steps, four other secured sets were identified for this printer: components 9, 10 and 11; components 12, 13 and 14; components 15, 16 and 17; and components 2, 3 and 4. Just for visualization purposes, the secured sets were rearranged, placing the secured set of components 2, 3 and 4 in the first row of the $SecM^2$, instead of the fifth row, and displacing the other rows down. This is why, in Tables 3 and 4, the set of components 6, 7 and 8 appears as Set 2, but the order is irrelevant in the application of the framework. In this manner the secured sets are defined for all printers.

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Table 6. Results summary for the printer HP Business Inkjet 2600. _

		X3(2)]		[Y3						Se	ecN	13					
Cluster 13	Cluster 12	Cluster 11	Cluster 10	Cluster 9	Cluster 8	Cluster 7	Cluster 6	Cluster 5	Cluster 4	Cluster 3	Cluster 2	Cluster 1		Yc3(2)		Y3(2)	Y3(1)	Y1(0)		Set 9	Set 8	Set 7	Set 6	Set 5	Set 4	Set 3	Set 2	Set 1		Comp. #
0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	ł	0	0	0		0	0	0	0	0	0	0	0	0	Carriage	
0	0	0	0	0	0	0	0	0	0	0	0	-		0	ľ	0	0	0		0	0	0	0	0	0	0	0	1	Plastic case K Ph	2
0	0	0	0	0	0	0	0	0	0	0	0	н		0	Ì	0	0	0		0	0	0	0	0	0	0	0	1	Printhead K	ω
0	0	0	0	0	0	0	0	0	0	0	н	0	1	0	Ì	0	0	0	İ	0	0	0	0	0	0	0	н	0	Plastic case C Ph	4
0	0	0	0	0	0	0	0	0	0	0	-	0	1	0	Ì	0	0	0		0	0	0	0	0	0	0	P	0	Printhead C	ъ
0	0	0	0	0	0	0	0	0	0	-	0	0	1	0	Ì	0	0	0		0	0	0	0	0	0	н	0	0	Plastic case M Ph	6
0	0	0	0	0	0	0	0	0	0	н	0	0	1	0	Ì	0	0	0		0	0	0	0	0	0	н	0	0	Printhead M	7
0	0	0	0	0	0	0	0	0	н	0	0	0	1	0	Ì	0	0	0		0	0	0	0	0	н	0	0	0	Plastic case Y Ph	∞
0	0	0	0	0	0	0	0	0		0	0	0	1	0	Ì	0	0	0		0	0	0	0	0	1	0	0	0	Printhead Y	9
0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	ľ	0	0	0		0	0	0	0	0	0	0	0	0	Carriage PCB	10
0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	ł	0	0	0		0	0	0	0	0	0	0	0	0	PCB Holder	11
0	0	0	0	÷	0	0	0	0	0	0	0	0	1	щ		1	ч	0		0	0	0	0	0	0	0	0	0	Belt attach	12
0	0	0	0	0	0	0	0	0	0	0	0	0		0		0	0	0		0	0	0	0	0	0	0	0	0	Secure Printheads Lever	13
0	0	0	0	0	0	0	0	0	0	0	0	0		0	ł	0	0	0		0	0	0	0	0	0	0	0	0	Flex cable	14
0	0	0	1	0	0	0	0	0	0	0	0	0		1		1	1	0		0	0	0	0	0	0	0	0	0	Carriage Belt	13
0	0	0	0	0	0	0	0	0	0	0	0	0		0		0	0	0		0	0	0	0	0	0	0	0	0	Slider	16
0	0	-	0	0	0	0	0	0	0	0	0	0		1		1	1	0		0	0	0	0	0	0	0	0	0	Carriage Motor	5
0		0	0	0	0	-	0	0	0	0	0	0		0		0	0	0		-	0	0	0	0	0	0	0	0	Ink Delivery System	L.
0	0	0	0	0	0	0	0	0	0	0	0	0		0	ł	0	0	0		0	0	0	0	0	0	0	0	0	Ink Supply Station Body	10
	0	0	0	0	0	0	0	0	0	0	0	0		0	ł	0	0	0		-	0	0	0	0	0	0	0	0	ISS PCB	22
		0	0	0	0	-	0	0	0	0	0	0		-	ł	0		0			0	0	-	0	0	0	0	0	Ink Pump	2
	0	0	0	0	0		0	-	0	0	0	0		0			1	0		0	0	0	0	-	0	0	0	0	Ink depository K	1 2
			0	0	-		0		-			0		-				0		E	-	0	-		-	0	-		Plastic case K	2 2
			0	0	0		0			0		0		-			-	0		E	0	0		-	0	0	0		Smart chin K	3 2
			0	0	0					0		0		-			-	0			0	0		0	0	0	0	0		4 2
-		-	-	0	-			0	-			0		-	ł	100	-	-		-		0	-	0	-	-	-	0	Plastic case C	5 2
	0		0	0	0		-	-	0	-	-	0		-			-	0		-	0	0	-	0	0	0	0	0	Flastic case C	6 2
0	0	0	0	0	0		-		0	0		0		0	ł	1	-	0			0	0		0	0	0	0	0	Smart chip C	7 2
0	0	0	0	0	0	1	0	0	0	0	0	0		0	ł	10	4	0		0	0	-	0	0	0	0	0	0		00
0	0	0	0	0	0		0	0	0	0	0	0		0			1	0		0	0	1	0	0	0	0	0	0	Plastic case M	ق
0	0	0	0	0	0	1	0	0	0	0	0	0		0	ł	1	1	0		-	0	P	0	0	0	0	0	0	Smart chip ivi	õ
0	0	0	0	0	-	0	0	0	0	0	0	0		0		1		0		0	-	0	0	0	0	0	0	0	Ink depository Y	1
0	0	0	0	0	1	0	0	0	0	0	0	0		0	-	1	1	0		-	P	0	0	0	0	0	0	0	Plastic case Y	22
0	0	0	0	0	P	0	0	0	0	0	0	0		0		1	1	0		0	-	0	0	0	0	0	0	0	Smart chip Y	8
0	0	0	0	0	0	0	0	0	0	0	0	0		0	ł	0	0	0		0	0	0	0	0	0	0	0	0	Serv. Station Body	4
0	0	0	0	0	0	0	0	0	0	0	0	0		0		0	0	0		0	0	0	0	0	0	0	0	0	Top Case	5
0	0	0	0	0	0	0	0	0	0	0	0	0		0		0	0	0		0	0	0	0	0	0	0	0	0	Wiper	36
-	0	0	0	0	0	0	0	0	0	0	0	0		4		E.	1	0		0	0	0	0	0	0	0	0	0	Sponge	37
0	0	0	0	0	0	0	0	0	0	0	0	0		0		0	0	0		0	0	0	0	0	0	0	0	0	Pallet	38
0	0	0	0	0	0	0	0	0	0	0	0	0		0	ļ	0	0	0		0	0	0	0	0	0	0	0	0	Shaft	39
0	0	0	0	0	0	0	0	0	0	0	0	0		0		0	0	0		0	0	0	0	0	0	0	0	0	Motor & Gears	40
0	0	0	0	0	T	1	1	1	0	0	0	0	Yn	n3(2	:)															



Figure 3. Functional structure and definition of the critical path for the printing subsystem.

In this work it is assumed that, usually, the critical functions in a product are related to the functions that require more research, development and resources from the company; therefore, identifying those functions tends to be a straightforward decision. The analysis performed to establish the set of critical functions for this illustrative example is an approximation to the company's perspective on the functional structure of the printing subsystems considered in this article.

4.3. Results

The results for the model of the HP Deskjet 1220C are shown in Table 5, for the HP Photosmart Plus are in Table 4, and for the HP Business Inkjet 2600 are in Table 6. The tables include (in addition to the *SecM*'s) the decision vectors at each iteration ($Y^k(t)$), the final modified decision vectors ($\hat{Y}^k(2)$), and the final chromosomes ($X^k(2)$) with the optimal clustering strategy.

The security matrices were generated for the critical functions selected in the *CrFun* vector, using the individual mapping from functions to components for each printer captured in the FCM's. Most of the secure sets are related to the ink cartridges and the printheads as expected; however, for the Business Inkjet 2600 the ninth set includes components 18 and 21 (the Ink Delivery System and the Ink Pump, respectively) which could have been easily overlooked if a manual selection had been done instead of the automated generation (see Table 6).

The overall process was completed in only two iterations of the master problem (component sharing). The decision vectors reflect the suitable candidates for sharing throughout the process without the restriction on security. These decision vectors are shown to illustrate how sensitive components could become suitable candidates for sharing if no consideration for security is taken. The final modified decision vector indicates the recommended set of components to share in each printer. When the security restrictions are involved in the optimization process, the decision vectors are modified and the constraint imposed with the *SecM* is followed strictly, which leads to a different component arrangement.

The restricted sets from the *SecM*'s appear preserved in the optimal clusters of each printer (see Tables 4, 5 and 6). One of the components to share is common for the three printers, the *Sponge*; some other components are shared only between two of the printers. The set of components to share between printers 2 (Photosmart Plus) and 3 (Business Inkjet 2600) include the *Carriage Belt*, the *Belt Attach*, and the *Carriage Motor*. These components are not sensitive and are not restricted by the sets defined. However, the ink cartridges are also candidates for sharing between these two printers, which are sensitive components. Sharing these elements is done in modules rather than the individual components; therefore, all the components in the restricted set should be candidates for sharing, which is the case for the ink cartridges.

Cluster 3 of printer 1 does not include any sensitive element or candidate for sharing, but corresponds to an optimal arrangement according to the connectivity of the system; thence, it is preserved by the algorithm.

When security restrictions were not considered (strictly following the framework in Rojas Arciniegas and Kim 2011), the results showed a more complicated module definition in the sense that some modules overlap to minimize mismatches of types I and II; however, when security is taken into account, the overlap is forbidden by the constraints in *SecM* or some modules break down leading to a simpler structure but greater mismatches of types I and II. The main difference was that the carriage of each printer was integrated with the components that constitute the cartridges or the printheads. (The results of this analysis following Rojas Arciniegas and Kim 2011 are not shown in this article owing to space limitations.) These results confirm how the security considerations affect the architecture of the product, and how the framework presented in this article, allows the designers to define sets of components that are known to be sensitive and therefore need to be enclosed in a module.

5. Conclusion and future work

The methodology proposed in this article takes component/module security into considerations for component sharing and optimal architecture decision making in a family of products.

This new security criterion affects significantly the outcome of the optimal component sharing decision making process in two ways. First, it preserves the restricted sets of components together while achieving optimal architecture. This may involve adding related components to the restricted sets to form a cluster, or even joining restricted sets into a bigger cluster, as long as it contributes to achieving optimal architecture. Second, it prevents the sensitive components from becoming candidates for sharing due to security concerns defined by the manufacturer. Instead, the sensitive components are allowed to be shared only as modules confined within the restricted set(s).

The illustration design example of an inkjet printer family shows how the security concern can be preserved while allowing component sharing across multiple products. This example validates the proposed framework by obtaining a different product architecture that considers both the security of the sensitive components and component sharing among the three printers. The results were compared against the application of a similar framework without the security considerations (although it is not shown owing to space limitations) and not only were they different but the results from the proposed framework reflect much more closely the actual implementation by HP and suggest which components could be shared.

It is not clear if the product architecture decisions that HP has made, at least for the three printers considered in this example, were indeed influenced by security considerations, but they certainly make sense. The cartridge of the entry level printer (which is also the most probable candidate for copying owing to high volume sales and high willingness from the user to buy the cheapest option) is the most complex type that they offer, enclosing the printhead as well as the ink reservoir under the same case, making it more difficult to copy. On the other hand, the

high-end printer offers simpler cartridges that only contain the ink and a separate cartridge for the printhead. This high-end type of printer is also intended for professional users, or at least users with very high concerns for quality and performance, who are less likely to buy third-party cartridges that may not meet the performance offered by the OEM.

Future work involves developing new representations with flexible degrees of security, which would capture security sensitive design decisions in a more comprehensive manner. These design decisions involve evaluating the tradeoffs between compromising the security of the system and gaining the advantages of sharing components in a family of products, for which there is no methodology established in the literature. Similarly, further research is necessary to expand this framework to design products or subsystems that comply with the anti-tamper requirements for products or components that are not strictly software or electronics for military applications.

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